

# A SHORT NOTE ON THE LACKING POLYNOMIAL OF THE COMPLETE BIPARTITE GRAPH

AMAL ALOFI AND MARK DUKES\*

ABSTRACT. We classify the stochastically recurrent states of the stochastic sandpile model on the complete bipartite graphs  $K_{2,n}$  and  $K_{m,2}$ . We use these characterizations to give formulae for the lacking polynomials on these graphs. Log-concavity of the sequence of coefficients of these two lacking polynomials is proven.

The stochastic sandpile model (SSM) is a variant of the Abelian sandpile model [2] that was introduced by Chan, Marckert and Selig [3] in 2013. In the SSM the classical sandpile toppling rule is replaced with an alternative rule whereby, on toppling an unstable vertex, a grain may (but does not have to) be sent to each neighbouring vertex. One consequence of this is that the toppling of a vertex does not necessarily result in it becoming stable. As part of their paper the authors introduced a notion of stochastically recurrent states as the analog of recurrent states for the classical model. They provided an expression for such states in terms of *orientations compatible with configurations*. Further research into this model can be found in the papers [6, 7].

Let  $G$  be a simple graph with a sink  $s$ . A *stable configuration* on  $G$  is an assignment of non-negative integers to each non-sink vertex such that the number of grains at a given vertex is less than its degree. First, we will recall a definition from Chan et al. [3] which explains what it means for an orientation to be compatible with a configuration.

**Definition 1** (Chan et al. [3]). Let  $c$  be the configuration on  $G$ . An orientation  $\mathcal{O}$  of  $G$  is an assignment of a direction to each of the edges of  $G$ . We say that configuration  $c$  is *compatible with  $\mathcal{O}$*  (and likewise  $\mathcal{O}$  is *compatible with  $c$* ) if for all non-sink vertices  $v_i$  in  $G$ ,

$$in_{\mathcal{O}}(v_i) \geq d(v_i) - c_i,$$

where  $in_{\mathcal{O}}(v_i)$  is the number of incoming edges to  $v_i$  w.r.t.  $\mathcal{O}$ .

We denote the set of stable configurations on  $G$  that are compatible with  $\mathcal{O}$  as  $comp(\mathcal{O})$ . Note that, in comparison to the paper [3], the inequality in Definition 1 is missing one on the right hand side. This is because of a subtle change to the model. In [3] they considered a vertex unstable if the number of grains at a vertex exceeds its degree, whereas we consider a vertex unstable if the number of grains at vertex is not less than the degree which is in line with the definition of the ASM in [4]. The two notions are equivalent.

**Theorem 2** (Chan et al. [3]). A stable configuration  $c$  on  $G$  is *stochastically recurrent* if and only if there exists an orientation  $\mathcal{O}$  on  $G$  such that  $c \in comp(\mathcal{O})$ . In other words,

$$Sto(K_{m,n}) = \bigcup_{\substack{\text{orientations } \mathcal{O} \\ \text{of } K_{m,n}}} comp(\mathcal{O})$$

where the union is taken over all orientations on  $G$ .

---

*Key words and phrases.* Stochastic sandpile model; classification problem; lacking polynomial; complete bipartite graph; log-concavity.

Chan et al. [3] also introduced the lacking polynomial of a graph to be the generating function counting the stochastically recurrent configurations according to the number of grains by which a given configuration differs from the maximally stable configuration.

**Definition 3** (Chan et al. [3]). The lacking polynomial  $L_G(x)$  is

$$L_G(x) := \sum_{c \in \text{Sto}(G)} x^{\ell(c)},$$

where

$$\ell(c) := \sum_{v \in V(G) \setminus \{s\}} l(v)$$

and  $l(v) := d(v) - c(v) - 1$  is the *lacking number* of vertex  $v$ .

In this note we consider the SSM on the complete bipartite graph  $K_{m,n}$  with vertex set  $\{v_0, v_1, \dots, v_{m+n-1}\}$ . We will treat  $v_0$  as the sink and this graph has edges connecting vertices in the sets  $\{v_0, v_1, v_2, \dots, v_{m-1}\}$  and  $\{v_m, v_{m+1}, \dots, v_{m+n-1}\}$ .

We characterise stochastically recurrent states on the graphs  $K_{2,n}$  and  $K_{m,2}$  and use these characterisations to give expressions for the lacking polynomials  $L_{2,n}(x)$  and  $L_{m,2}(x)$  on those graphs. We also prove that the sequence of coefficients of both  $L_{2,n}(x)$  and  $L_{m,2}(x)$  are log-concave. This note is motivated by Alofi and Dukes [1] that considers rectangular tableaux representations of recurrent states of the Abelian sandpile model on the complete bipartite graph, and transformations upon them.

Theorem 2 allows us to write an expression for stochastically recurrent states on the complete bipartite graph  $K_{m,n}$ :

**Proposition 4.** The set of stochastically recurrent states of the SSM on  $K_{m,n}$  is

$$\text{Sto}(K_{m,n}) = \bigcup_{\substack{\text{orientations } \mathcal{O} \\ \text{of } K_{m,n}}} \{(c_1, \dots, c_{m+n-1}) : \text{out}_{\mathcal{O}}(v_i) \leq c_i < d(v_i), \forall 1 \leq i \leq m+n-1\}.$$

*Proof.* Let  $c = (c_1, c_2, \dots, c_{m+n-1})$  be a stable configuration on  $K_{m,n}$ . Suppose it is compatible with an orientation  $\mathcal{O}$  where  $c_i < d(v_i)$  for all  $1 \leq i \leq m+n-1$ . According to the Definition 1 we must have:

$$\text{in}_{\mathcal{O}}(v_i) \geq d(v_i) - c_i.$$

This means  $c_i \geq d(v_i) - \text{in}_{\mathcal{O}}(v_i) = \text{out}_{\mathcal{O}}(v_i)$ . When we combine the application of Definition 1 with Theorem 2 we find

$$\text{Sto}(K_{m,n}) = \bigcup_{\substack{\text{orientations } \mathcal{O} \\ \text{of } K_{m,n}}} \{(c_1, \dots, c_{m+n-1}) : \text{out}_{\mathcal{O}}(v_i) \leq c_i < d(v_i), \forall 1 \leq i \leq m+n-1\}.$$

□

**Example 5.** Consider the graph  $K_{2,2}$ . To determine the stochastically recurrent states for each orientation  $\mathcal{O}$  of graph  $K_{2,2}$  first we find  $\text{out}_{\mathcal{O}}(v_i)$  for all  $1 \leq i \leq 3$ , and then when we apply Prop. 4. See Table 1 for the derivations for each orientation. It follows that

$$\text{Sto}(K_{2,2}) = \{(0, 1, 1), (1, 0, 1), (1, 1, 0), (1, 1, 1)\}.$$

We can provide crude lower and upper bounds on the number of stochastically recurrent configurations by making use of the fact [3, Prop. 2.3] that  $\text{Sto}(K_{m,n})$  properly contains the set of classically recurrent states, and such states are in 1-1 correspondence with the number of spanning tree of the underlying graph. Fieldler and Sedlacek [5] showed the number of spanning trees of the complete bipartite graph  $K_{m,n}$  is  $n^{m-1}m^{n-1}$ . Thus

	$out_{\mathcal{O}}(v_1) = 1, out_{\mathcal{O}}(v_2) = 1, out_{\mathcal{O}}(v_3) = 1.$ $\implies 1 \leq c_1 < 2, 1 \leq c_2 < 2, \text{ and } 1 \leq c_3 < 2.$ The contribution to $\text{Sto}(K_{2,2})$ is $\{(1, 1, 1)\}.$
	$out_{\mathcal{O}}(v_1) = 1, out_{\mathcal{O}}(v_2) = 1, out_{\mathcal{O}}(v_3) = 1.$ $\implies 1 \leq c_1 < 2, 1 \leq c_2 < 2, \text{ and } 1 \leq c_3 < 2.$ The contribution to $\text{Sto}(K_{2,2})$ is $\{(1, 1, 1)\}.$
	$out_{\mathcal{O}}(v_1) = 1, out_{\mathcal{O}}(v_2) = 0, out_{\mathcal{O}}(v_3) = 1.$ $\implies 1 \leq c_1 < 2, 0 \leq c_2 < 2, \text{ and } 1 \leq c_3 < 2.$ The contribution to $\text{Sto}(K_{2,2})$ is $\{(1, 0, 1), (1, 1, 1)\}.$
	$out_{\mathcal{O}}(v_1) = 1, out_{\mathcal{O}}(v_2) = 1, out_{\mathcal{O}}(v_3) = 0.$ $\implies 1 \leq c_1 < 2, 1 \leq c_2 < 2, \text{ and } 0 \leq c_3 < 2.$ The contribution to $\text{Sto}(K_{2,2})$ is $\{(1, 1, 0), (1, 1, 1)\}.$
	$out_{\mathcal{O}}(v_1) = 0, out_{\mathcal{O}}(v_2) = 1, out_{\mathcal{O}}(v_3) = 1.$ $\implies 0 \leq c_1 < 2, 1 \leq c_2 < 2, \text{ and } 1 \leq c_3 < 2.$ The contribution to $\text{Sto}(K_{2,2})$ is $\{(0, 1, 1), (1, 1, 1)\}.$

TABLE 1. Checking all orientations of  $K_{2,2}$ .

the number of stochastically recurrent configurations on  $K_{m,n}$  is at least  $n^{m-1}m^{n-1}$ . A trivial upper bound is achieved by noting the stochastically recurrent states are stable configurations, of which there are  $n^{m-1}m^n$  many. Thus

$$n^{m-1}m^{n-1} \leq |\text{Sto}(K_{m,n})| \leq n^{m-1}m^n. \quad (1)$$

Note that the upper bound differs from the lower bound only by a factor of  $m$ .

**Question 6.** Can it be determined whether or not the number of stochastically recurrent states dominates the set of stable states? I.e. can it be decided

$$|\text{Sto}(K_{m,n})| \leq \frac{n^{m-1}m^n}{2}?$$

The set of stable configurations on  $K_{2,n}$  is

$$\{(c_1, c_2, \dots, c_{n+1}) : 0 \leq c_1 < n \text{ and } c_i \in \{0, 1\}, \forall 2 \leq i \leq n+1\}.$$

**Theorem 7.** The lacking polynomial of the graph  $K_{2,n}$  is

$$L_{2,n}(x) = \sum_{k=0}^{n-1} \sum_{i=0}^k \binom{n}{i} x^k.$$

*Proof.* Definition 3 gives

$$L_{2,n}(x) = \sum_{c \in \text{Sto}(K_{2,n})} x^{\ell(c)}$$

where

$$\ell(c) := \sum_{v_i \in V(K_{2,n}) \setminus \{v_0\}} l(v_i) \quad \text{and} \quad l(v_i) := d(v_i) - c(v_i) - 1.$$

In order to calculate this sum we need to have an explicit characterization of the set  $\text{Sto}(K_{2,n})$ . Using Proposition 4 we know

$$\text{Sto}(K_{2,n}) = \{c = (c_1, \dots, c_{n+1}) : \exists \text{ an orientation } \mathcal{O} \text{ of } K_{2,n} \text{ compatible with } c\}.$$

Let us suppose  $c = (c_1, \dots, c_{n+1})$  is a member of  $\text{Sto}(K_{2,n})$ . Then we must have that there exists an orientation  $\mathcal{O}$  of  $K_{2,n}$  that is compatible with  $c$ . This means

$$n - \text{in}_{\mathcal{O}}(v_1) = \text{out}_{\mathcal{O}}(v_1) \leq c_1 < n$$

and

$$2 - \text{in}_{\mathcal{O}}(v_j) = \text{out}_{\mathcal{O}}(v_j) \leq c_j < 2 \text{ for all } 2 \leq j \leq n + 1.$$

Let us now see under what conditions one can construct an orientation  $\mathcal{O}$  on  $K_{2,n}$  that is compatible with a given  $c$ . Let  $X$  be the set of indices  $j$  in  $[2, n + 1]$  for which  $c_j = 0$ , and where we use the notation  $[a, b] := \{a, a + 1, \dots, b\}$ . For any  $j$  in  $X$  the number of outgoing edges at vertex  $v_j$  is zero ( $\text{out}_{\mathcal{O}}(v_j) = 0$ ), because we know that  $\text{out}_{\mathcal{O}}(v_j) \leq c_j < 2$ , so if  $c_j = 0$  then  $\text{out}_{\mathcal{O}}(v_j) = 0$ . Therefore, for all  $j$  in  $X$  there is one outgoing edge from  $v_1$  to  $v_j$ , and hence vertex  $v_j$  has one incoming edge from  $v_1$ . So the number of outgoing edges from  $v_1$  is greater than or equal to the number of elements in  $X$ .

For any  $j$  in  $[2, n + 1] \setminus X$ , we know that there is at most one outgoing edge from  $v_j$  as  $\text{out}_{\mathcal{O}}(v_j) \leq c_j < 2$ . So if  $c_j = 1$  then  $\text{out}_{\mathcal{O}}(v_j) \leq 1$ . With these considerations in mind, we can construct an orientation  $\mathcal{O}$  on  $K_{2,n}$  that is compatible with a stable configuration  $c$ :

- (i) If  $c_j = 1$  with  $2 \leq j \leq n + 1$  then there is at most one outgoing edge from  $v_j$ .
- (ii) If  $c_j = 0$  and  $2 \leq j \leq n + 1$  then there are no outgoing edges from  $v_j$ .
- (iii) The outgoing edges from vertex  $v_1$  is greater than or equal the the number of vertices  $v_j$  when  $c_j = 0$  for all  $2 \leq j \leq n + 1$ . We know that the  $\text{out}_{\mathcal{O}}(v_1)$  is less than or equal to  $c_1$  and greater than or equal to the number of vertices  $v_j$  when  $c_j = 0$  for all  $2 \leq j \leq n + 1$ . Therefore  $c_1$  is greater than or equal to the number of vertices  $v_j$  when  $c_j = 0$  for all  $2 \leq j \leq n + 1$ .

There are no other restrictions that forbid us from constructing such an orientation  $\mathcal{O}$ . Therefore we can write down the following self-contained expression for the set  $\text{Sto}(K_{2,n})$  that does not depend on an orientation  $\mathcal{O}$ :

$$\text{Sto}(K_{2,n}) = \{(c_1, c_2, \dots, c_{n+1}) : c_2, \dots, c_{n+1} \in \{0, 1\} \text{ and } c_1 \geq |j \in [2, n + 1] : c_j = 0|\}.$$

Now that we have an explicit expression for the set  $\text{Sto}(K_{2,n})$ , we can use this expression to calculate the lacking polynomial.

Let  $c$  be in  $\text{Sto}(K_{2,n})$ . Let  $l_1 = n - c_1 - 1$  be the lacking number at vertex  $v_1$ , and let  $l_j = 1 - c_j$  be the lacking number at vertices  $v_2, \dots, v_{n+1}$  for all  $j \in [2, n + 1]$ , so  $l_j = 0$  when  $c_j = 1$  and  $l_j = 1$  when  $c_j = 0$ . Let  $i$  be the number of vertices  $v_j$  with  $j \in [2, n + 1]$  and  $c_j = 0$  for which  $l_j = 1$ . The remaining vertices have  $l_j = 0$ . Then  $x^{\ell(c)}$  factors as  $x^{\ell(c)} = x^i x^{l_1}$ . Since  $c_1 \geq |j \in [2, n + 1] : c_j = 0|$ , we conclude that  $c_1 \geq i$ , therefore the lacking number at  $v_1$  will be between 0 and  $n - 1 - i$ . Since there are  $\binom{n}{i}$  combinations of  $i$  vertices with lacking number 1 among  $\{v_2, \dots, v_{n+1}\}$ , we obtain

$$L_{2,n}(x) = \sum_{i=0}^n \binom{n}{i} x^i \sum_{l_1=0}^{n-1-i} x^{l_1}.$$

Now setting  $k = i + l_1$  we have

$$L_{2,n}(x) = \sum_{k=0}^{n-1} \sum_{i=0}^k \binom{n}{i} x^k.$$

□

A configuration  $c = (c_1, c_2, \dots, c_{m+1})$  on  $K_{m,2}$  is stable precisely when  $c_i \in \{0, 1\}$  for all  $i \in \{1, \dots, m-1\}$  and  $0 \leq c_j < m$  for all  $j \in \{m, m+1\}$ .

**Theorem 8.** The lacking polynomial for the graph  $K_{m,2}$  is

$$L_{m,2}(x) = \sum_{k=0}^{m-1} S(m-1, k)x^k$$

where

$$S(m-1, k) = \sum_{q=0}^k \sum_{r=0}^q \binom{m-1}{r} \text{ for all } 0 \leq k \leq m-1.$$

*Proof.* The lacking polynomial  $L_{m,2}(x)$ , given in Definition 3, is

$$L_{m,2}(x) = \sum_{c \in \text{Sto}(K_{m,2})} x^{\ell(c)}$$

where

$$\ell(c) = \sum_{v_i \in V(K_{m,2}) \setminus \{v_0\}} l(v_i), \quad \text{and} \quad l(v_i) = d(v_i) - c(v_i) - 1.$$

To calculate this we require an explicit characterization of the set  $\text{Sto}(K_{m,2})$ . From Proposition 4 we have

$$\text{Sto}(K_{m,2}) = \{c = (c_1, \dots, c_{m+1}) : \exists \text{ orientation } \mathcal{O} \text{ of } K_{m,2} \text{ compatible with } c\}.$$

Suppose  $c = (c_1, \dots, c_{m+1})$  is a member of  $\text{Sto}(K_{m,2})$ . Then we must have that there exists an orientation  $\mathcal{O}$  of  $K_{m,2}$  that is compatible with  $c$ . This means that

$$n - in_{\mathcal{O}}(v_i) = out_{\mathcal{O}}(v_i) \leq c_i < 2 \text{ for all } i \in \{1, \dots, m-1\}$$

and

$$2 - in_{\mathcal{O}}(v_j) = out_{\mathcal{O}}(v_j) \leq c_j < m \text{ for all } j \in \{m, m+1\}.$$

Let us see under what conditions one can construct an orientation  $\mathcal{O}$  on  $K_{m,2}$  that is compatible with a given  $c$ . Let  $X$  be the set of indices  $i$  in  $[1, m-1]$  for which  $c_i = 0$ . For any  $i \in X$  the number of outgoing edges at vertex  $v_i$  is zero ( $out_{\mathcal{O}}(v_i) = 0$ ) since  $out_{\mathcal{O}}(v_i) \leq c_i < 2$ . So if  $c_i = 0$  then  $out_{\mathcal{O}}(v_i) = 0$ .

Therefore, for all  $i$  in  $X$  there are is one outgoing edge from each of  $v_m$  and  $v_{m+1}$  to  $v_i$ . Moreover, vertex  $v_i$  has one incoming edge from each of  $v_m$  and  $v_{m+1}$ . So the number of outgoing edges from  $v_m$  is greater than or equal to the number of elements in  $X$ . Also, the number of outgoing edges from  $v_{m+1}$  is greater than or equal to the number of elements in  $X$ . Therefore  $c_m \geq |X|$  and  $c_{m+1} \geq |X|$ .

For any  $i \in [1, m-1] \setminus X$ , we know that there is at most one outgoing edge from  $v_i$  since  $out_{\mathcal{O}}(v_j) \leq c_j < 2$ , so if  $c_j = 1$  then  $out_{\mathcal{O}}(v_j) \leq 1$ . So the total number of outgoing edges from  $v_m$  and  $v_{m+1}$  to  $v_i$  must at least equal  $m-1 - |X|$ . Therefore we can then construct such an orientation  $\mathcal{O}$  of  $K_{m,2}$  that is compatible with a stable configuration  $c$  precisely when  $c_m + c_{m+1} \geq 2|X| + m - 1 - |X| = m - 1 + |X|$ . There are no other restrictions that forbid us from constructing such an orientation  $\mathcal{O}$ . Therefore we can write down the following self-contained expression for the set  $\text{Sto}(K_{m,2})$  that does not depend on an orientation  $\mathcal{O}$ :

$$\text{Sto}(K_{m,2}) = \{(c_1, c_2, \dots, c_{m+1}) : c_1, \dots, c_{m-1} \in \{0, 1\} \text{ and } c_m + c_{m+1} \geq m - 1 + |X|\}.$$

Now that we have an explicit expression for the set  $\text{Sto}(K_{m,2})$ , we can use this expression to calculate the lacking polynomial. Let  $c$  be in  $\text{Sto}(K_{m,2})$ . Let  $l_i = 1 - c_i$  be the lacking

number of vertex  $v_i$  for all  $i \in [1, m-1]$ , so that  $l_i = 0$  when  $c_i = 1$  and  $l_i = 1$  when  $c_i = 0$ . For the vertices  $v_m$  and  $v_{m+1}$  the lacking numbers are  $l_m = m-1-c_m$  and  $l_{m+1} = m-1-c_{m+1}$ , respectively. We factor  $x^{\ell(c)} = x^{|X|}x^{l_m+l_{m+1}}$  and can now write  $c_m + c_{m+1} \geq m-1 + |X|$  in terms of lacking numbers as:

$$m-1-l_m+m-1-l_{m+1} \geq m-1+|X|$$

which equivalent to

$$m-1-|X| \geq l_m+l_{m+1}.$$

Now suppose that  $r = l_m + l_{m+1}$ , then  $r$  ranges in between 0 and  $m-1-|X|$ . For each choice of  $r$  we have exactly  $r+1$  choices for  $(l_m, l_{m+1})$ . Now let  $j = |X|$  be the number of vertices  $v_i$  with  $i \in [1, m-1]$  and  $c_i = 0$  for which then  $l_i = 1$ . The remaining vertices will have  $l_i = 0$ . So there are  $\binom{m-1}{j}$  combinations of  $j$  vertices with lacking number 1 among  $\{v_1, \dots, v_{m-1}\}$ , and we obtain

$$L_{m,2}(x) = \sum_{j=0}^{m-1} \binom{m-1}{j} x^j \sum_{r=0}^{m-1-j} (r+1)x^r = \sum_{j=0}^{m-1} \sum_{r=0}^{m-1-j} \binom{m-1}{j} (r+1)x^{r+j}.$$

Set  $k = j + r$ . Then  $k$  runs from 0 to  $m-1$  so that  $j$  can run from 0 to  $k$  and, with  $r = k - j$ , we obtain the sum

$$L_{m,2}(x) = \sum_{k=0}^{m-1} \sum_{j=0}^k \binom{m-1}{j} (k-j+1)x^k.$$

This is equal to

$$L_{m,2}(x) = \sum_{k=0}^{m-1} S(m-1, k)x^k$$

where

$$S(m-1, k) = \sum_{q=0}^k \sum_{r=0}^q \binom{m-1}{r} \text{ for all } 0 \leq k \leq m-1. \quad \square$$

A sequence  $a_0, a_1, \dots, a_n$  of non-negative real numbers is said to be *logarithmically concave*, or *log-concave*, if for all  $0 < k < n$ ,  $a_k^2 - a_{k-1}a_{k+1} \geq 0$ .

**Lemma 9.** Suppose  $(x_0, x_1, \dots, x_n)$  is a log-concave sequence. Then the sequence  $(z_0, \dots, z_n)$  of partial sums defined by

$$z_k = \sum_{i=0}^k x_i$$

is also log-concave.

*Proof.* Wang and Yeh [8] proved that if sequences  $(x_k)$  and  $(y_k)$  are log-concave, then the sequence  $(z_k)$  where  $z_k$  is the ordinary convolution

$$z_k = \sum_{i=0}^k x_i y_{n-i}$$

is log-concave. The sequence  $(y_0, \dots, y_k) = (1, \dots, 1)$  is trivially log-concave. The statement of the lemma follows since it is the special case with all  $y_j$ 's replaced with 1s.  $\square$

**Theorem 10.**

- (i) The sequence of coefficients of the lacking polynomial  $L_{2,n}(x)$  is log-concave.
- (ii) The sequence of coefficients of the lacking polynomial  $L_{m,2}(x)$  is log-concave.

*Proof.* (i) From Theorem 7 we have

$$L_{2,n}(x) = \sum_{k=0}^{n-1} T(n, k)x^k \quad \text{where} \quad T(n, k) := \sum_{i=0}^k \binom{n}{i}.$$

It is well-known that a sequence of binomial coefficients  $\left(\binom{n}{k}\right)_{k=0,1,2,\dots,n}$  is log-concave. By Lemma 9, the sequence  $(T(n, k))_{k=0,\dots,n}$  is log-concave. That is, for any  $0 < k \leq n-1$

$$(T(n, k))^2 - T(n, k-1)T(n, k+1) \geq 0$$

Therefore the sequence of coefficients of the lacking polynomial  $L_{2,n}(x)$  is log-concave.

(ii) From Theorem 8 we have

$$L_{m,2}(x) = \sum_{k=0}^{m-1} S(m-1, k)x^k$$

where

$$S(m-1, k) := \sum_{q=0}^k \sum_{r=0}^q \binom{m-1}{r} \quad \text{for all } 0 < k \leq m-1.$$

By Lemma 9 the sequence  $R(m-1, q) := \left(\sum_{r=0}^q \binom{m-1}{r}\right)_{q=0,\dots,m-1}$  is log-concave. Then, again by an application of Lemma 9, the sequence  $(S(m-1, k) = \sum_{q=0}^k R(m-1, q))_{k=0}^{m-1}$  is also log-concave. Therefore the sequence of coefficients of the lacking polynomial  $L_{m,2}(x)$  is log-concave.  $\square$

The two results in Theorem 10 suggest that log-concavity might be a property of these lacking polynomials for the general complete bipartite graph.

**Conjecture 11.** The sequence of coefficients of the lacking polynomial  $L_{m,n}(x)$  for the complete bipartite graph  $K_{m,n}$  is log-concave for all  $m, n \geq 2$ .

## REFERENCES

- [1] A. Alofi and M. Dukes. Parallelogram polyominoes and rectangular EW-tableaux: correspondences through the sandpile model. *Enumer. Comb. Appl.* **1** (2021), no. 1, Article S2R8.
- [2] D. Dhar. Theoretical studies of self-organized criticality. *Physica A Stat. Mech. Appl.* **369** (2006), no. 1, 29–70.
- [3] Y. Chan, J-F. Marckert, and T. Selig. A natural stochastic extension of the sandpile model on a graph. *J. Comb. Theory Ser. A* **120** (2013), no. 7, 1913–1928.
- [4] M. Dukes and Y. Le Borgne. Parallelogram polyominoes, the sandpile model on a complete bipartite graph, and a  $q, t$ -Narayana polynomial. *J. Comb. Theory Ser. A* **120** (2013), no. 4, 816–842.
- [5] M. Fieldler and J. Sedlacek. O  $W$ -basích orientovaných grafu. *Časopis. Pěst. Mat.* **83** (1958), 214–225.
- [6] T. Selig. The stochastic sandpile model on complete graphs. *arXiv:2209.07301*.
- [7] T. Selig and H. Zhu. Asymptotics of the single-source stochastic sandpile model. *arXiv:2208.10202*.
- [8] Y. Wang and Y-N. Yeh. Log-concavity and LC-positivity. *J. Comb. Theory Ser. A* **114** (2007), no. 2, 195–210.

SCHOOL OF MATHEMATICS AND STATISTICS, UNIVERSITY COLLEGE DUBLIN, BELFIELD, DUBLIN 4, IRELAND.

*Email address:* amal.alofi@ucdconnect.ie

*Email address:* mark.dukes@ucd.ie