

GRAPHS AND NETWORKS (MATH20150)

Problem sheet 4

1. This molecule is a tree with $m + n$ vertices. Therefore it has $m + n - 1$ edges. Since every C has 4 bonds and every H has one bond, the sum of the degrees is $4m + n$. By the degree sum formula we get that the number of edges is $(4m + n)/2$, i.e. $4m + n = 2(m + n - 1)$, so $n = 2m + 2$.
2. $A + I_n$ is the adjacency matrix of the graph G' obtained from G by adding a loop at each vertex.

“ \Rightarrow ” Consider the entry (i, j) of $(A + I_n)^{n-1}$. To show that it is non-zero we must show that there is a path of length $n - 1$ in G' from v_i to v_j . Since G is connected, there is a path P in G from v_i to v_j , and this path has length $k \leq n - 1$ (since there are only n vertices in G). To get a path of length $n - 1$ in G' from v_i to v_j , follow first the path P , then loop $n - 1 - k$ times at v_j (using the loop at v_j).

“ \Leftarrow ” Let v_i and v_j be vertices in G . Since $(A + I_n)^{n-1}$ has no zero entries, there is a path of length $n - 1$ in G' from v_i to v_j . Obviously we can remove all the loops from this path, and we obtain a path in G from v_i to v_j (we do not know its length).

3. Every time we remove an edge from a cycle the resulting graph is still connected (seen in class). So the graph T is connected. Since it contains no cycles, it is a tree. It still has n vertices (we did not remove any vertex from G), so T has $n - 1$ edges. Therefore we removed $m - n + 1$ edges.
4. “ \Rightarrow ” If x and y are in the same component of G , since this component is a tree, we know that there is exactly one path from x to y . If x and y are in different components of G , there is no path from x to y .

“ \Leftarrow ” Let H be a component of G . We show that H is a tree. Let $x, y \in H$. Since H is connected there is at least one path from x to y , and by hypothesis we know that there is exactly one path from x to y . So H is a tree, and G is a forest (since every component is a tree).